

A review of quadrature-based bounds of the error norms in CG

G rard MEURANT

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To set the notation, A symmetric of order n

Lanczos algorithm

input A, v

$$\beta_0 = 0, v_0 = 0$$

$$v_1 = v / \|v\|$$

for $k = 1, \dots$ **do**

$$w = Av_k - \beta_{k-1}v_{k-1}$$

$$\alpha_k = v_k^T w$$

$$w = w - \alpha_k v_k$$

$$\beta_k = \|w\|$$

$$v_{k+1} = w / \beta_k$$

end for

It generates tridiagonal matrices T_k , $k = 1, \dots, n$ with coefficients α_i, β_i

A symmetric positive definite, $Ax = b$

CG algorithm

input A, b, x_0

$$r_0 = b - Ax_0$$

$$p_0 = r_0$$

for $k = 1, \dots$ until convergence **do**

$$\gamma_{k-1} = \frac{r_{k-1}^T r_{k-1}}{p_{k-1}^T A p_{k-1}}$$

$$x_k = x_{k-1} + \gamma_{k-1} p_{k-1}$$

$$r_k = r_{k-1} - \gamma_{k-1} A p_{k-1}$$

$$\delta_k = \frac{r_k^T r_k}{r_{k-1}^T r_{k-1}}$$

$$p_k = r_k + \delta_k p_{k-1}$$

end for

Relations between CG and Lanczos

$$T_k = L_k D_k L_k^T$$

with

$$L_k \equiv \begin{pmatrix} 1 & & & & \\ \sqrt{\delta_1} & \ddots & & & \\ & \ddots & \ddots & & \\ & & \ddots & \ddots & \\ & & & \sqrt{\delta_{k-1}} & 1 \end{pmatrix}, \quad D_k \equiv \begin{pmatrix} \gamma_0^{-1} & & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ & & & & \gamma_{k-1}^{-1} \end{pmatrix}$$

$$\beta_k = \frac{\sqrt{\delta_k}}{\gamma_{k-1}}, \quad \alpha_k = \frac{1}{\gamma_{k-1}} + \frac{\delta_{k-1}}{\gamma_{k-2}}, \quad \delta_0 = 0, \quad \gamma_{-1} = 1$$

$$v_{j+1} = (-1)^j \frac{r_j}{\|r_j\|}, \quad j = 0, \dots, k$$

M.R. Hestenes and E. Stiefel, *Methods of conjugate gradients for solving linear systems*, J. Nat. Bur. Standards, v 49 n 6 (1952), pp. 409–436

Hestenes and Stiefel “error function” was the A -norm of the error

$$\|\varepsilon_k\|_A \equiv (\varepsilon_k^T A \varepsilon_k)^{1/2}$$

$$\|\varepsilon_k\|_A = \|x - x_k\|_A = \min_{y \in x_0 + \mathcal{K}_k(A, r_0)} \|x - y\|_A$$

Note that

$$\|\varepsilon_k\|_A^2 = r_k^T A^{-1} r_k, \quad \|\varepsilon_k\|^2 = r_k^T A^{-2} r_k$$

H-S also noticed the connection of CG with orthogonal polynomials and Riemann-Stieltjes integrals

Let

$$A = U\Lambda U^T, \quad UU^T = U^T U = I$$

with $\Lambda = \text{diag}(\lambda_1, \dots, \lambda_n)$, w be a given unit norm vector,

$$\omega_i \equiv (w, u_i)^2 \quad \text{so that} \quad \sum_{i=1}^n \omega_i = 1$$

and the stepwise constant distribution function

$$\omega(\lambda) \equiv \begin{cases} 0 & \text{for } \lambda < \lambda_1, \\ \sum_{j=1}^i \omega_j & \text{for } \lambda_i \leq \lambda < \lambda_{i+1}, \quad 1 \leq i \leq n-1, \\ 1 & \text{for } \lambda_n \leq \lambda \end{cases}$$

$$\int_{\mu}^{\eta} f(\lambda) d\omega(\lambda) = \sum_{i=1}^n \omega_i f(\lambda_i) = w^T f(A) w$$

Choosing $w = r_k / \|r_k\|$ and $f(\lambda) = 1/\lambda$, it is clear that $r_k^T A^{-1} r_k$ can be written as a Riemann-Stieltjes integral

This result was used by [Gene Golub](#) and his collaborators to compute approximations of quadratic forms $u^T f(A)v$ with several different applications by using Gauss quadrature rules

[G. Dahlquist, S.C. Eisenstat, and G.H. Golub](#), Bounds for the error of linear systems of equations using the theory of moments, J. Math. Anal. Appl., 37, pp. 151-166, 1972

[G. Dahlquist, G.H. Golub, and S.G. Nash](#) Bounds for the error in linear systems, in R. Hettich, ed., Proceedings of the workshop on semi-infinite programming, pp. 154-172, Berlin, 1978 Springer

B. Fischer and G.H. Golub, On the error computation for polynomial based iteration methods, in A. Greenbaum and M. Luskin, eds., Recent advances in iterative methods, pp. 59-67, 1994, Springer

G.H. Golub and Z. Strakoš, Estimates in quadratic formulas, Numer. Algorithms, 8(2), pp. 241-268, 1994

G.H. Golub and G. Meurant, Matrices, moments and quadrature, in D.F. Griffiths and G.A. Watson, eds., Numerical Analysis 1993, volume 303 of Pitman Research Notes in Mathematics, pp. 105-156, Longman Sci. Tech., 1994

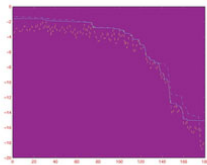
G.H. Golub and G. Meurant, Matrices, moments and quadrature II or how to compute the norm of the error in iterative methods, BIT, 37(3), pp. 687-705, 1997

G. Meurant, The computation of bounds for the norm of the error in the conjugate gradient algorithm, Numer. Algorithms, 16, pp. 77-87, 1997

The computation of $u^T f(A)v$ and applications are summarized in the book published by Princeton University Press in 2010

Princeton Series in APPLIED MATHEMATICS

Matrices, Moments and Quadrature with Applications



Gene H. Golub
G rard Meurant

For CG we used the relation

$$\|\varepsilon_k\|_A^2 = \|r_0\|^2 [(T_n^{-1}e_1, e_1) - (T_k^{-1}e_1, e_1)]$$

It was shown by Z. Strakoš and P. Tichý that this relation holds in finite precision arithmetic up to a small perturbation term

It can also be written as

$$\|\varepsilon_0\|_A^2 = \sum_{j=0}^{k-1} \gamma_j \|r_j\|^2 + \|\varepsilon_k\|_A^2$$

or

$$(T_n^{-1})_{1,1} = (T_k^{-1})_{1,1} + \mathcal{R}_k^{(G)}[\lambda^{-1}]$$

$$(T_n^{-1})_{1,1} = \frac{\|\varepsilon_0\|_A^2}{\|r_0\|^2} = \int_{\mu}^{\eta} \lambda^{-1} d\omega(\lambda)$$

$(T_k^{-1})_{1,1}$ is the Gauss quadrature approximation of the integral and the remainder is

$$\mathcal{R}_k^{(G)}[\lambda^{-1}] = \frac{\|\varepsilon_k\|_A^2}{\|r_0\|^2}$$

The nodes of the Gauss quadrature rule are the eigenvalues of T_k

Since we don't know $\|\varepsilon_0\|_A^2$, we use

$$\|\varepsilon_{k-d}\|_A^2 - \|\varepsilon_k\|_A^2 = \sum_{j=k-d}^{k-1} \gamma_j \|r_j\|^2$$

where d is a positive integer smaller than k

The right-hand side is a lower bound of the A -norm squared at iteration $k - d$

To obtain an upper bound of the A -norm we use a Gauss-Radau quadrature rule with a fixed node $\mu < \lambda_1$

$$\|\varepsilon_0\|_A^2 = \|r_0\|^2 (\widehat{T}_{k+1}^{-1})_{1,1} + \widehat{\mathcal{R}}_{k+1}[\lambda^{-1}]$$

Subtracting

$$\|\varepsilon_0\|_A^2 = \|r_0\|^2 (T_{k-d}^{-1})_{1,1} + \|\varepsilon_{k-d}\|_A^2$$

we obtain

$$\|\varepsilon_{k-d}\|_A^2 = \|r_0\|^2 [(\widehat{T}_{k+1}^{-1})_{1,1} - (T_{k-d}^{-1})_{1,1}] + \widehat{\mathcal{R}}_{k+1}[\lambda^{-1}]$$

The difference in the right-hand side can be written as

$$(\widehat{T}_{k+1}^{-1})_{1,1} - (T_{k-d}^{-1})_{1,1} = (\widehat{T}_{k+1}^{-1})_{1,1} - (T_k^{-1})_{1,1} + Q_{k-d,d}$$

with the Gauss lower bound

$$Q_{k-d,d} = \sum_{j=k-d}^{k-1} \gamma_j \|r_j\|^2$$

We have to find $\alpha_{k+1}^{(\mu)}$ such that μ is an eigenvalue of the extended tridiagonal matrix

$$\widehat{T}_{k+1}^{(\mu)} = \begin{pmatrix} \alpha_1 & \beta_1 & & & & & \\ \beta_1 & \ddots & \ddots & & & & \\ & \ddots & \ddots & & & & \\ & & & \beta_{k-1} & & & \\ & & & \beta_{k-1} & \alpha_k & \beta_k & \\ & & & & \beta_k & \alpha_{k+1}^{(\mu)} & \end{pmatrix}$$

where the α_j 's and β_j 's are the Lanczos coefficients
It is known that

$$\alpha_{k+1}^{(\mu)} = \mu + \xi_k^{(\mu)}$$

where $\xi_k^{(\mu)}$ is the last component of the solution of

$$(T_k - \mu I)\xi^{(\mu)} = \beta_k^2 e_k$$

This is computed with the LDL^T factorization of $T_k - \mu I$

Then, we can use the [Sherman-Morrison](#) for the difference
 $(\hat{T}_{k+1}^{-1})_{1,1} - (T_k^{-1})_{1,1}$

This gives the CGQL algorithm in

[G.H. Golub and G. Meurant, 1997](#)

CG coeffs \rightarrow Lanczos coeffs \rightarrow Gauss-Radau upper bound

Can we compute the upper bound directly from the CG coefficients?

We look for a coefficient $\gamma_k^{(\mu)}$ such that

$$T_{k+1}^{(\mu)} = L_{k+1} \begin{pmatrix} D_k & \\ & (\gamma_k^{(\mu)})^{-1} \end{pmatrix} L_{k+1}^T$$

such that μ is an eigenvalue

This problem was solved in

G. Meurant and P. Tichý, On computing quadrature-based bounds for the A -norm of the error in conjugate gradients, Numer. Algorithms, 62(2), pp. 163-191, 2013

$$\gamma_{j+1}^{(\mu)} = \frac{\gamma_j^{(\mu)} - \gamma_j}{\mu(\gamma_j^{(\mu)} - \gamma_j) + \delta_{j+1}}, \quad \gamma_0^{(\mu)} = \frac{1}{\mu}$$

This leads to the CGQ algorithm

It was also proved that

$$\gamma_k \|r_k\|^2 < \|\varepsilon_k\|_A^2 < \gamma_k^{(\mu)} \|r_k\|^2 < \left(\frac{\phi_k}{\mu}\right) \|r_k\|^2$$

with $\phi_k = \|r_k\|^2 / \|p_k\|^2$ and

$$\phi_k = \frac{\phi_{k-1}}{\phi_{k-1} + \delta_k}, \quad \phi_0 = 1$$

input A, b, x_0, μ, η, d

$$r_0 = b - Ax_0, p_0 = r_0, g_0^{(\mu)} = \frac{\|r_0\|^2}{\mu}, g_0^{(\eta)} = \frac{\|r_0\|^2}{\eta}$$

for $k = 1, \dots$ until convergence do

$$\gamma_{k-1} = \frac{r_{k-1}^T r_{k-1}}{p_{k-1}^T A p_{k-1}}$$

$$x_k = x_{k-1} + \gamma_{k-1} p_{k-1}, r_k = r_{k-1} - \gamma_{k-1} A p_{k-1}$$

$$\delta_k = \frac{r_k^T r_k}{r_{k-1}^T r_{k-1}}$$

$$p_k = r_k + \delta_k p_{k-1}$$

$$g_{k-1} = \gamma_{k-1} \|r_{k-1}\|^2$$

$$\Delta_{k-1}^{(\mu)} = g_{k-1}^{(\mu)} - g_{k-1}, \quad g_k^{(\mu)} = \frac{\|r_k\|^2 \Delta_{k-1}^{(\mu)}}{\mu \Delta_{k-1}^{(\mu)} + \|r_k\|^2}$$

$$\Delta_{k-1}^{(\eta)} = g_{k-1}^{(\eta)} - g_{k-1}, \quad g_k^{(\eta)} = \frac{\|r_k\|^2 \Delta_{k-1}^{(\eta)}}{\eta \Delta_{k-1}^{(\eta)} + \|r_k\|^2}$$

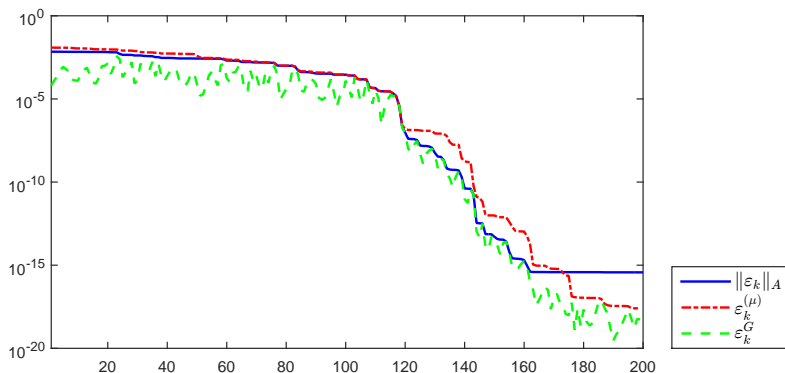
$$Q_{k-d,d} = \sum_{j=k-d}^{k-1} g_j, \quad \varepsilon_{k-d}^G = \sqrt{Q_{k-d,d}}$$

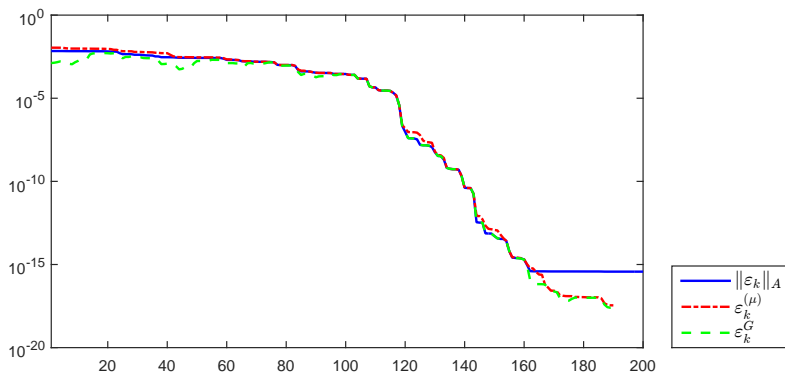
$$\varepsilon_{k-d}^{(\mu)} = \sqrt{Q_{k-d,d} + g_k^{(\mu)}}, \quad \varepsilon_{k-d}^{(\eta)} = \sqrt{Q_{k-d,d} + g_k^{(\eta)}}$$

end for

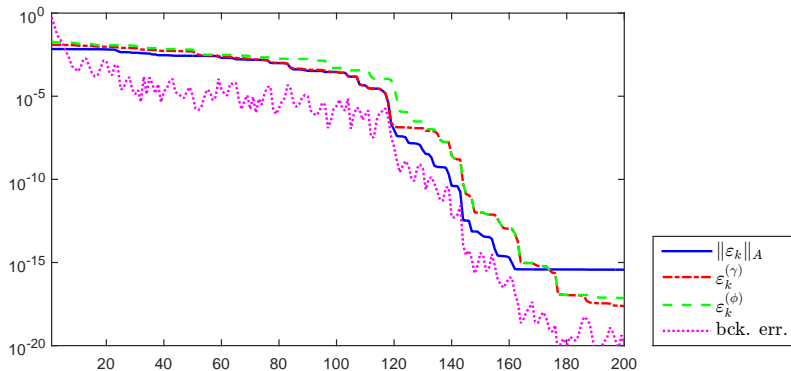
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$\mu = (1 - 10^{-8})\lambda_1 = 3417.267528494$, whence
 $\lambda_1 = 3417.267562666$, $d = 1$



Same μ , $d = 10$ 

Same μ , other bounds



How to choose d ?

An algorithm for choosing the delay for the Gauss lower bound is described in

G. Meurant, J. Papež, and P. Tichý, Accurate error estimation in CG, Numer. Algorithms, 88(3), pp. 1337-1359, 2021

to obtain

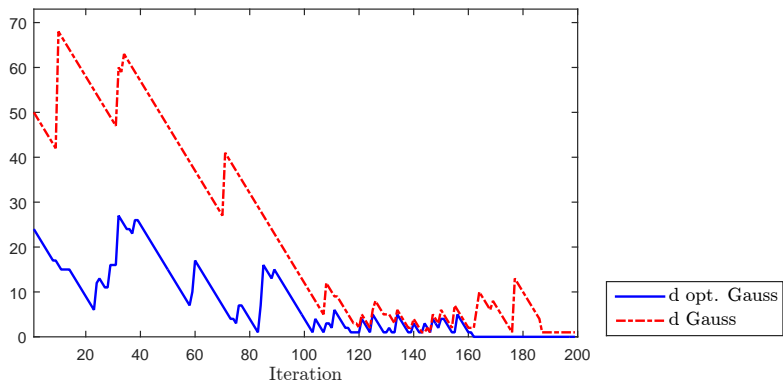
$$\frac{\|\varepsilon_{k-d}\|_A^2 - Q_{k-d,d}}{\|\varepsilon_{k-d}\|_A^2} \leq \tau$$

with

$$Q_{k-d,d} = \sum_{j=k-d}^{k-1} \gamma_j \|r_j\|^2$$

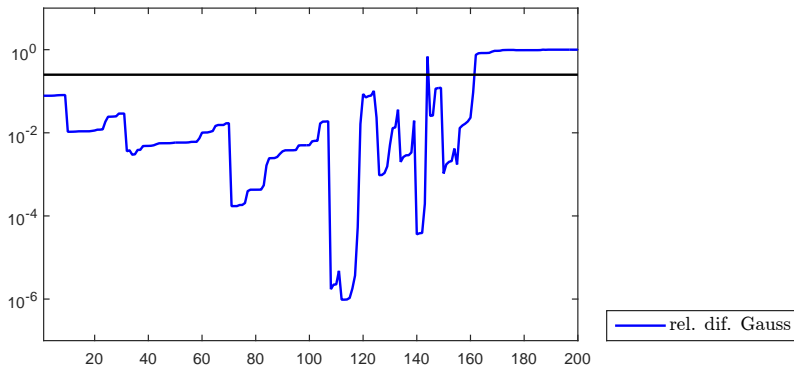
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$$\tau = 0.25$$



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$$\tau = 0.25$$



The Gauss-Radau problem

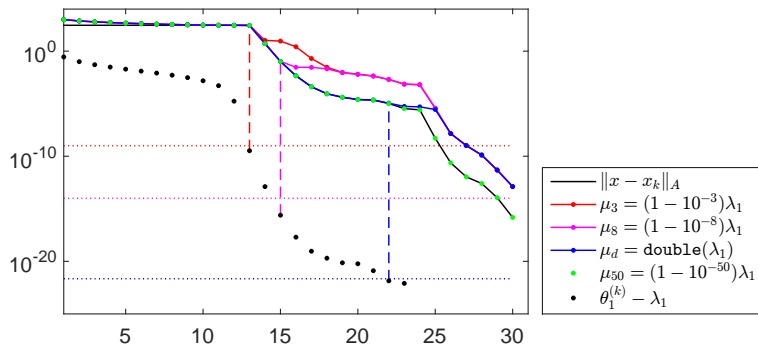
To study why the Gauss-Radau bound is not tight even when μ is close to λ_1 , we set up a model problem of order 30 and we run CG in extended precision with `digits=128`

G. Meurant and P. Tichý, The behaviour of the Gauss-Radau upper bound of the A -norm of the error for the conjugate gradient algorithm, in preparation

The problem does not come from rounding errors

Model problem

$\theta_1^{(k)}$ is the smallest Ritz value



The problem starts when $\theta_1^{(k)} - \lambda_1 < \lambda_1 - \mu$

We have a theoretical explanation for that phenomenon

One remedy is to increase the delay for Gauss-Radau

We are currently working on an adaptative algorithm to do so

Other papers on error estimation:

[C. Brezinski](#), Error estimates for the solution of linear systems, SIAM J. Sci. Comput., 21, pp. 764-781, 1999

[D. Calvetti](#), [S. Morigi](#), [L. Reichel](#), and [F. Sgallari](#), Computable error bounds and estimates for the conjugate gradient method, Numer. Algorithms, 25, pp. 79-88, 2000

[Z. Strakoš](#) and [P. Tichý](#), Error estimation in preconditioned conjugate gradients, BIT Numerical Mathematics, 45, pp. 789-817, 2005

[A. Frommer](#), [K. Kahl](#), [T. Lippert](#), and [H. Rittich](#), 2-norm error bounds and estimates for Lanczos approximations to linear systems and rational matrix functions, SIAM J. Matrix Anal. Appl., 34(3), pp. 1046-1065, 2013

[R. Estrin](#), [D. Orban](#), and [M.A. Saunders](#), Euclidean-norm error bounds for SYMMLQ and CG, SIAM J. Matrix Anal. Appl., 40(1), pp. 235-253, 2019

[E. Hallman](#) Sharp 2-norm error bounds for LSQR and the conjugate gradient method, SIAM J. Matrix Anal. Appl., 41(3), pp. 1183-1207, 2020